

# Sharp Bounds for Spectral Radius of Graphs Presented by $K$ -Neighbour of the Vertices

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## Abstract

Let  $G = (V, E)$  be a simple connected graph with  $V(G) = \{v_1, v_2, \dots, v_n\}$  and degree sequence  $d_1, d_2, \dots, d_n$ . Denote

$$t_k(i) = \sum_{(v_i, v_j) \in E} t_{k-1}(j), m_k(i) = \frac{t_{k+1}(i)}{t_k(i)}, \text{ where } k \text{ is a positive}$$

integer number and  $v_i \in V(G)$  and note that  $t_0(i) = d_i$ . Let  $\rho(G)$  be the largest eigenvalue of adjacent matrix of  $G$ . In this paper, we present sharp upper and lower bounds of  $\rho(G)$  in terms of  $m_k(i)$  (see theorem (2.1)). From which, we can obtain some known results, and our result is better than other results in some case.

**Key words:** Spectral radius; Bound;  $k$ -neighbour

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## INTRODUCTION

Let  $G = (V, E)$  be a connected graph without loops and multiedges and vertex set  $V = \{v_1, v_2, \dots, v_n\}$ . The degree  $d_i$  of a vertex  $v_i$  in the graph  $G$  is defined to be the number of edges in  $G$  adjacent to  $v_i$ . For  $v_i \in V(G)$ ,  $N(v_i)$  denotes the neighbors of  $v_i$ . The 2-degree of  $v_i$  (Brualdi & Hoffman, 1985) is the sum of the degrees of the vertices adjacent to  $v_i$  and denoted by  $t_i$ , and the average-degree of  $v_i$  is

$$m_i = \frac{t_i}{d_i}. \text{ Here we define}$$

$$t_k(i) = \sum_{(v_i, v_j) \in E} t_{k-1}(j), m_k(i) = \frac{t_{k+1}(i)}{t_k(i)},$$

where  $k$  is a positive integer number. Note that

$$t_0(i) = d_i, m_0(i) = m_i.$$

Let  $A(G) = (a_{ij})$ ,  $a_{ij} = 1$  if  $(v_i, v_j) \in E$  and  $a_{ij} = 0$  otherwise be the adjacency matrix of  $G$ . It follows immediately that if  $G$  is a simple graph, then  $A(G)$  is a symmetric  $(0, 1)$  matrix in which every diagonal entry is zero. Since  $A(G)$  is real and symmetric, its eigenvalues are real. The spectral radius of  $G$ , denoted by  $\rho(G)$ , is the largest eigenvalue of  $A(G)$ . Note that if  $G$  is connected, then  $A(G)$  is irreducible, and so by the Perron-Frobenius theory of non-negative matrices,  $\rho(G)$  has multiplicity one and there exists a unique positive unit eigenvector (also called Perron-eigenvector) corresponding to  $\rho(G)$ .

Up to now, many bounds for  $\rho(G)$  were given. For example, Kinkar Ch.Das and Pawan Kumar (Das & Kumar, 2004) gave a bound of spectral radius for graphs:

$$\min\{\sqrt{m_i m_j} : ij \in E\} \leq \rho(G) \leq \max\{\sqrt{m_i m_j} : ij \in E\}, \quad (1)$$

where  $m_i$  is the average degree of  $v_i$ . Moreover, the equality holds if and only if  $G$  is either a graph with all the vertices of equal average degree or a bipartite graph with vertices of same set having equal average degree.

In this paper, we will generalize the Kinkar Ch.Das and Pawan Kumar's bound and obtain the upper and lower bounds on  $\rho(G)$  in terms of  $m_k(i)$ . From which, we can obtain some known results (for example (1)). We will give an example to show that our result is better than the bound (1) in some case.

Now we introduce some lemmas which will be used later on.

**Lemma 1.1** (Horn & Johnson, 1985). *Let  $A$  be a nonnegative matrix of order  $n$ .  $R_i$  be the  $i$ th row sum of  $A$ . Then*

$$\min\{R_i : 1 \leq i \leq n\} \leq \rho(A) \leq \max\{R_i : 1 \leq i \leq n\}.$$

*If  $A$  is irreducible, then each equality holds if and only*

if  $R_1=R_2=\dots=R_n$ .

**Lemma 1.2** Let  $G$  be a bipartite graph with bipartition  $V = U \cup W$  and  $m_k(i)=\alpha$  for  $v_i \in U$ ,  $m_k(j)=\beta$  for  $v_j \in W$ , Then,  $\rho(G) = \sqrt{\alpha\beta}$ .

*Proof.* Let  $A$  be the adjacency matrix of  $G$ . Obviously,  $\rho(G)$  is the spectral radius of the matrix  $M=K^{-1}(D^{-1}AD)K$  too, where  $D=\text{diag } t_k(1), t_k(2), \dots, t_k(n)$ ,  $K = \text{diag}(k_1, k_2, k_n)$ ,  $k_i = \sqrt{\alpha}$  when  $v_i \in U$  and  $k_i = \sqrt{\beta}$  when  $v_i \in W$ ,  $1 \leq i \leq n$ . Then  $(i, j)$ th element of the matrix  $M$  is equal to

$$\begin{cases} \sqrt{\frac{\beta}{\alpha}} \frac{t_k(j)}{t_k(i)} & \text{if } (v_i, v_j \in E), v_i \in U; \\ \sqrt{\frac{\alpha}{\beta}} \frac{t_k(j)}{t_k(i)} & \text{if } (v_i, v_j \in E), v_i \in W; \\ 0 & \text{otherwise.} \end{cases}$$

So each row sum of the matrix  $M$  is equal to  $\sqrt{\alpha\beta}$ . Thus, by Lemma 1.1, we have  $\rho(G) = \sqrt{\alpha\beta}$ .

## THE BOUNDS OF SPECTRAL RADIUS

**Theorem 2.1** Let  $G$  be a connected graph. Then

$$\min \{ \sqrt{m_k(i)m_k(j)} : (v_i, v_j) \in E \} \leq \rho(G) \leq \max \{ \sqrt{m_k(i)m_k(j)} : (v_i, v_j) \in E \}. \quad (2)$$

Moreover, either of equality holds for a particular value of  $k$  if and only if  $m_k(1)=m_k(2)=\dots=m_k(n)$  or  $G$  is a bipartite graph with the partition  $\{v_1, v_2, \dots, v_{n_1}\} \cup \{v_{n_1+1}, \dots, v_n\}$  and  $m_k(1)=m_k(2)=\dots=m_k(n_1)$ ,  $m_k(n_1+1)=\dots=m_k(n)$ .

*Proof.* It is easy to see that the proof of lower bound is similar as the upper bound, so we only give the proof of upper bound. Let  $D = \text{diag } t_k(1), t_k(2), \dots, t_k(n)$ . Obviously,  $D^{-1}AD$  and  $A$  have the same spectral radius. Let  $X=(x_1, x_2, \dots, x_n)^T$  be an eigenvector of  $D^{-1}AD$  corresponding to the spectral radius  $\rho(G)$ . Let one eigencomponent (say  $x_i$ ) be equal to 1 and the other eigencomponents be less than or equal to 1, that is,  $x_1 = 1$ , and  $0 < x_k \leq 1$  for all  $k$ . Let  $x_2 = \max \{ x_k : (v_1, v_k) \in E \} \geq \max \{ x_k : (v_i, v_k) \in E \}$  when  $x_i = 1$ .

Now the  $(i, j)$ th element of  $D^{-1}AD$  is  $\frac{t_k(j)}{t_k(i)}$ , if  $(v_i, v_j) \in E$ ,

and 0 otherwise.

We have

$$D^{-1}ADX = \rho(G)X. \quad (3)$$

From the first equation of (3), we have

$$\begin{aligned} \rho x_1 &= \sum_{(v_1, v_j) \in E} \frac{t_k(j)x_j}{t_k(1)}, \\ \rho &\leq m_k(1)x_2. \end{aligned} \quad (4)$$

From the second equation of (3), we have

$$\begin{aligned} \rho x_2 &= \sum_{(v_2, v_j) \in E} \frac{t_k(j)x_j}{t_k(2)}, \\ \rho x_2 &\leq m_k(2). \end{aligned} \quad (5)$$

From (4) and (5), we get

$$\rho^2 \leq m_k(1)m_k(2).$$

Therefore,

$$\rho^2 \leq \sqrt{m_k(1)m_k(2)}.$$

Hence,

$$\rho(G) \leq \max \{ \sqrt{m_k(i)m_k(j)} : (v_i, v_j) \in E \}.$$

Now suppose that the equality in (2) holds. Then all inequalities in the above argument must be equalities. In particular, we have from (4) that  $x_j=x_2$  for all  $k$ ,  $(v_i, v_j) \in E$ , also from (5) that  $x_j=x_1=1$  for all  $k$ ,  $(v_2, v_j) \in E$ . Now we distinguish two cases below:

Case (i):  $x_2=1$ . Let  $V_1 = \{k : x_k=1\}$ . If  $V_1 \neq V$ , there exist vertices  $r, p \in V_1, q \notin V_1$  such that  $(v_r, v_p) \in E$ , and  $(v_p, v_q) \in E$ . since  $G$  is connected, so  $x_r=x_p=x_q=1$ .

From the  $r$ -th equation of (3), we have

$$\rho x_r = \sum_{(v_r, v_j) \in E} \frac{t_k(j)x_j}{t_k(r)} \leq m_k(r).$$

From the  $p$ -th equation of (3), we have

$$\rho x_p = \sum_{(v_p, v_j) \in E} \frac{t_k(j)x_j}{t_k(p)} < m_k(p).$$

So we have  $\rho(G) < \sqrt{m_k(r)m_k(p)}$  which contradicts that the equality holds in (2). Thus  $V_1=V$  and  $m_k(1)=m_k(2)=\dots=m_k(n)=\rho$ .

Case (ii):  $x_2 < 1$ . We have  $x_j=1, v_j \in N(v_2)$  and  $v_j=x_2, v_j \in N(v_1)$ . Let  $U = \{k : x_k=1\}$  and  $W = \{k : x_k=x_2\}$ , so  $N(v_1) \subseteq W$ , and  $N(v_2) \subseteq U$ . Further, for any vertex  $v_r \in N(N(v_1))$ , there exists a vertex  $v_p \in N(v_1)$ , such that  $(v_1, v_p) \in E$  and  $(v_p, v_r) \in E$ , therefore  $x_p=x_2$ . From the  $p$ -th equation of (3), we have

$$\rho x_p = \sum_{(v_p, v_j) \in E} \frac{t_k(j)x_j}{t_k(p)} \leq m_k(p),$$

Using (4), we get

$$\rho^2 \leq m_k(1)m_k(p),$$

since we have

$$\rho(G) = \max \{ \sqrt{m_k(i)m_k(j)} : (v_i, v_j) \in E \} \geq \sqrt{m_k(1)m_k(p)},$$

so  $\rho(G) < \sqrt{m_k(r)m_k(p)}$ , which shows that  $x_r=1$ . hence  $N(N(v_1)) \subseteq U$ . By a similar argument, we can show that  $N(N(v_1)) \subseteq W$ . Continuing the procedure, it is easy to see, since  $G$  is connected, that  $V = U \cup W$  and that the subgraphs induced by  $U$  and  $W$ , respectively, are empty

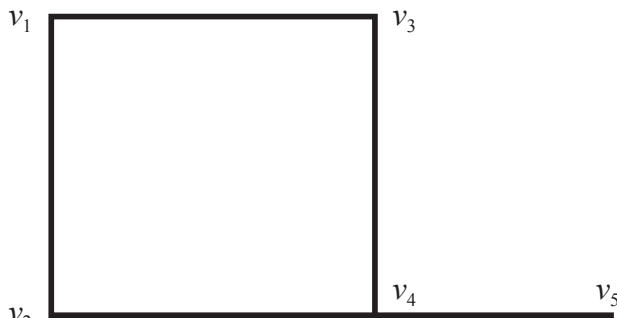
graphs. Hence  $G$  is bipartite and  $m_k(i)$  are the same for  $v_i \in U$ ,  $m_k(j)$  are the same for  $v_j \in W$ .

Conversely, if  $G$  is a graph with  $m_k(1)=m_k(2)=\dots=m_k(n)$ , then the equality in (2) is satisfied. Let  $G$  be a bipartite graph with bipartition  $V = U \cup W$  and  $m_k(i)=\alpha$  for  $v_i \in U$ ,  $m_k(j)=\beta$  for  $v_j \in W$ . Then, by lemma 1.2,

$$\rho(G) = \sqrt{\alpha\beta} = \max \{ \sqrt{m_k(i)m_k(j)} : (v_i, v_j) \in E \},$$

we complete the proof.

**Note 2.2** If  $k = 0$ , then the inequality (2) is the Kinkar Ch.Das and Pawan Kumar's bound (1). Here we give an example to show that (2) is better than the Kinkar Ch.Das and Pawan Kumar's bound in some case. Let  $G$  be a graph shown in Figure 1. Then the bound (2) is  $2.52 \leq \rho(G) \leq 2.67$  when  $k = 1$ , and Kinkar Ch.Das and Pawan Kumar's bound is  $2.49 \leq \rho(G) \leq 2.83$ . Thus in that case, (2) is better than the Kinkar Ch.Das and Pawan Kumar's bound.



**Figure 1**  
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Consequently, from (2) we have the following results.

**Corollary 2.3** Let  $G$  be a simple connected graph.

Then

$$\min \{ m_k(i) : i \in V \} \leq \rho(G) \leq \max \{ m_k(i) : i \in V \}. \quad (6)$$

Moreover, equality holds for a particular value of  $k$  if and only if  $m_k(1)=m_k(2)=\dots=m_k(n)$ .

**Note 2.4** If  $k = 0$ , then the inequality (6) is the Favaron et.al.,  $s$  bound (Favaron, Maheo, & Sacle, 1993, p.193).

**Corollary 2.5** (a) Let  $G$  be a graph with  $m_k(v)=p$ . For each  $v \in V(G)$ , then  $\rho(G)=p$ .

(b) Let  $G$  be a bipartite graph with the bipartition  $(X, Y)$ . If  $m_k(v)=p_x$  for each  $v \in X$  and  $m_k(v)=p_y$ . For each  $v \in Y$ , then  $\rho(G) = \sqrt{p_x p_y}$ .

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